

# Development and Verification of a Flight Stack for a High-Altitude Glider in Ada/SPARK 2014\*

Martin Becker, Emanuel Regnath, and Samarjit Chakraborty  
Chair of Real-Time Computer Systems, Technical University of Munich  
80333 Munich, Germany

**Abstract.** SPARK 2014 is a modern programming language and a new state-of-the-art tool set for development and verification of high-integrity software. In this paper, we explore the capabilities and limitations of its latest version in the context of building a flight stack for a high-altitude unmanned glider. Towards that, we deliberately applied static analysis early and continuously during implementation, to give verification the possibility to steer the software design. In this process we have identified several limitations and pitfalls of software design and verification in SPARK, for which we give workarounds and protective actions to avoid them. Finally, we give design recommendations that have proven effective for verification, and summarize our experiences with this new language.

**Keywords:** Ada/SPARK, formal verification, limitations, rules

## 1 Introduction

The system under consideration is a novel kind of weather balloon which is actively controlled, and thus requires verification to ensure it is working properly in public airspace. As any normal weather balloon, the system climbs up to the stratosphere (beyond an altitude of 10 km), while logging weather data such as temperature, pressure, NO<sub>2</sub>-levels and so on. Eventually the balloon bursts, and the sensors would be falling back to the ground with a parachute, drifting away with prevailing wind conditions. However, our system is different from this point onwards: the sensors are placed in a light-weight glider aircraft which is attached to the balloon. At a defined target altitude, the glider separates itself from the balloon, stabilizes its attitude and performs a controlled descent back to the take-off location, thus, bringing the sensors back home. In this paper we focus on the development and formal verification of the glider's onboard software.

The requirements for such a system are challenging already because of the extreme environmental conditions; temperatures range from 30 °C down to -50 °C, winds may exceed 100 kph, and GPS devices may yield vastly different output in

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\*The source code for this project is available at [github.com/tum-ei-rics/StratoX](https://github.com/tum-ei-rics/StratoX). This is a pre-print. The final publication will be available at Springer via [http://dx.doi.org/10.1007/978-3-319-50000-0\\_10](http://dx.doi.org/10.1007/978-3-319-50000-0_10) in *Computer Safety, Reliability and Security, 36th International Conference SAFECOMP 2017*, S. Tonetta and E. Schoitsch and F. Bitsch (Eds.), Trento, Italy, 2017.

those altitudes due to decreasing precision and the wind conditions. The combination of those extreme values is likely to trigger corner cases in the software, and thus should be covered by means of extensive testing or by analysis.

We use this opportunity of a safety-critical, yet hardly testable system to explore the new state-of-the-art verification tools of Ada/SPARK 2014 [7], especially to identify limitations, pitfalls and applicability in practice. To experience this new SPARK release to its full extent, we applied a *co-verification* approach. That is, we did not perform verification on a finished product, but instead in parallel to the software development (the specific strategy is not of relevance for this paper, but only the effect that this enabled us to identify code features that pose challenges in verification, and find workarounds for them). The implementation could therefore be shaped by verification needs. Moreover, since the high-altitude glider was a research project, we allowed ourselves to modify the initial software design to ease verification when needed.

## 2 Verification in Ada/SPARK

SPARK 2014 is a major redesign of the original SPARK language, which was intended for formal verification. SPARK 2014 now adopts Ada 2012 syntax, and covers a large subset of Ada. As a result, the GNAT Ada compiler can build an executable from SPARK 2014 source code, and even compile a program which mixes both languages. Compared to Ada, the most important exclusions are pointers (called *access*), aliasing and allocators, as well as a ban of exception handling. As a consequence, SPARK programs first and foremost must be shown to be free of run-time exceptions (called *AoRTE* - absence of run-time errors), which constitutes the main verification task.

The SPARK language – for the rest of this paper we refer to SPARK 2014 simply as SPARK – is built on functional contracts and data flow contracts. Subprograms (procedures and functions) can be annotated with pre- and post-conditions, as well as with data dependencies. GNATprove, the (only) static analyzer for SPARK 2014, aims to prove subprograms in a modular way, by analyzing each of them individually. The effects of callees are summarized by their post-condition when the calling subprogram is analyzed, and the precondition of the callee is imposing a proof obligation on the caller, i.e., the need to verify that the caller respects the callee’s precondition. Further proof obligations arise from each language-defined check that is executed on the target, such as overflow checks, index checks, and so on. If all proofs are successful, then the program is working according to its contracts and no exceptions will be raised during execution, i.e., AoRTE is established.

Internally, GNATprove [7] builds on the Why3 platform [6], which performs deductive verification on the proof obligations to generate verification conditions (VCs), and then passes them to a theorem solver of user’s choice, e.g., *cvc4*, *alt-ergo* or *z3*. Note that there exists also a tool for abstract interpretation, which is, however, not discussed here.

## 2.1 The GNAT Dimensionality Checking System

We also want to introduce a feature that is not part of the SPARK language itself, but an implementation-defined extension of the GNAT compiler, and thus available for SPARK programs. Since Ada 2012, the GNAT compiler offers a dimension system for numeric types through implementation-defined aspects [9]. The dimension system can consist of up to seven base dimensions, and physical quantities are declared as subtypes, annotated with the exponents of each dimension. Expressions using such variables are statically analyzed by the compiler for their dimensional consistency. Furthermore, the dimensioned variables contribute to readability and documentation of the code. Inconsistencies such as the following are found (dividend and divisor are switched in the calculation of rate):

```
1 | angle : Angle_Type := 20.0 * Degree;  
2 | dt    : Time_Type  := 100.0 * Milli * Second;  
3 | rate  : Angular_Velocity_Type := dt / angle; -- compiler error
```

Note that scaling prefixes like `milli` can be used, and that common conversions, such as between Degree and Radian in line 1, can be governed in a similar way.

In our project, we specified a unit system with the dimensions *length*, *mass*, *time*, *temperature*, *current*, and *angle*. Adding angle as dimension provides better protection against assignments of dimensionless types, as proposed in [11].

## 3 Initial System Design & Verification Goals

**Target Hardware.** We have chosen the “Pixhawk” autopilot [8]. It comprises two ARM processors; one Cortex-M4F (STM32F427) acting as flight control computer, and one Cortex-M3 co-processor handling the servo outputs. We implemented our flight stack on the Cortex-M4 from the ground up, thus completely replacing the original PX4/NuttX firmware that is installed when shipped.

### Board Support, Hardware Abstraction Layer & Run-Time System.

We are hiding the specific target from the application layer by means of a board support package (not to be confused with an Ada package). This package contains an hardware abstraction layer (HAL) and a run-time system (RTS). The RTS is implementing basic functionality such as tasking and memory management. The HAL is our extension of AdaCore’s Drivers Library [1], and the RTS is our port of the Ada RTS for the STM32F409 target. Specifically, we have ported the Ravenscar Small Footprint variant [3], which restricts Ada’s and SPARK’s tasking facilities to a deterministic and analyzable subset, but meanwhile forbids exception handling, which anyway is not permitted in SPARK.

**Separating Tasks by Criticality** has been one goal, since multi-threading is supported in SPARK. In particular, 1. termination of low-critical tasks must not cause termination of high-critical tasks, 2. higher-criticality tasks must not be blocked by lower-critical tasks and, 3. adverse effects such as deadlocks, priority inversion and race condition must not occur. We partitioned our glider software into two tasks (further concurrency arises from interrupt service routines):

1. The *Flight-Critical Task* includes all execution flows required to keep the glider in a controlled and navigating flight, thus including sensor reads and actuator writes. It is time-critical for control reasons. High-criticality.
2. The *Mission-Critical Task* includes all execution flows that are of relevance for recording and logging of weather data to an SD card. Low-priority task, only allowed to run when the flight-critical task is idle. Low-criticality.

The latter task requires localization data from the former one, to annotate the recorded weather data before writing it to the SD card. Additionally, it takes over the role of a flight logger, saving data from the flight-critical task that might be of interest for a post-flight analysis. The interface between these two tasks would therefore be a protected object with a message queue that must be able to hold different types of messages.

**Verification Goals.** First and foremost, AoRTE shall be established for all SPARK parts, since exceptions would result in task termination. Additionally, the application shall make use of as many contracts and checks as possible, and perform all of its computations using dimension-checked types. Last but not least, a few functional high-level requirements related to the homing functionality have been encoded in contracts. Overall, the focus of verification was the application, not the BSP. The BSP has been written in SPARK only as far as necessary to support proofs in the application. The rationale was that the RTS was assumed to be well tested, and the HAL was expected to be hardly verifiable due to direct hardware access involving pointers and restricted types.

## 4 Problems and Workarounds

In this section, we describe the perils and difficulties that we identified during verification of SPARK programs. We use the following nomenclature:

- **False Positive** denotes a failing check (failed VC) in static analysis which would not fail in any execution on the target, i.e., a false alarm.
- **False Negative** denotes a successful check (discharged VC) in static analysis which would fail in at least one execution on the target, i.e., a missed failure.

### 4.1 How to Miss Errors

There are a few situations in which static analysis can miss run-time exceptions, which in a SPARK program inevitably ends in abnormal program termination. Before we show these unwanted situations, we have to point out one important property of a deductive verification approach: Proofs build on each other. Consider the following example (results of static analysis given in comments):

```

1 | a := X / Z; -- medium: division check might fail
2 | b := Y / Z; -- info: division check proved

```

The analyzer reports that the check in line 2 cannot fail, although it suffers from the same defect as line 1. However, when the run-time check at line 1 fails, then

line 2 cannot be reached with the offending value of `z`, therefore line 2 is not a False Negative, unless exceptions have been wrongfully disabled.

**Mistake 1: Suppressing False Positives.** When a developer comes to the conclusion that the analyzer has generated a False Positive (e.g., due to insufficient knowledge on something that is relevant for a proof), then it might be justified to suppress the failing property. However, we experienced cases where this has generated False Negatives which were hiding (critical) failures. Consider the following code related to the GPS:

```

1 function toInt32 (b : Byte_Array) return Int_32 with Pre => b'Length = 4;
2 procedure Read_From_Device (d : out Byte_Array) is begin
3   d := (others => 0); -- False Positive
4   pragma Annotate (GNATprove, False_Positive, "length check might fail",
5     ...);
6 end Read_From_Device;
7
8 procedure Poll_GPS is
9   buf      : Byte_Array(0..91) := (others => 0);
10  alt_mm   : Int_32;
11 begin
12   Read_From_Device (buf);
13   alt_mm := toInt32(buf(60..64)); -- False Negative, guaranteed exception
14 end Poll_GPS;

```

Static analysis found that the initialization of the array `a` in line 3 could fail, but this is not possible in this context, and thus a False Positive<sup>1</sup>. The developer was therefore suppressing this warning with an annotation pragma. However, because proofs build on each other, a severe defect in line 12 was missed. The array slice has an off-by-one error which *guaranteed* failing the precondition check of `toInt32`. The reason for this False Negative is that everything after the initialization of `a` became virtually *unreachable* and that all following VCs consequently have been discharged. In general, a False Positive may exclude some or all execution paths for its following statements, and thus hide (critical) failure. We therefore recommend to avoid suppressing False Positives, and either leave them visible for the developer as warning signs, or even better, rewrite the code in a prover-friendly manner following the tips in Section 5.1.

**Mistake 2: Inconsistent Contracts.** Function contracts act as barriers for propagating proof results (besides inlined functions), that is, the result of a VC in one subprogram cannot affect the result of another in a different subprogram. However, these barriers can be broken when function contracts are inconsistent, producing False Negatives by our definition. One way to obtain inconsistent contracts, is writing a postcondition which itself contains a failing VC (line 2):

```

1 function f1 (X : Integer) return Integer
2   with Post => f1'Result = X + 1 is -- overflow check might fail
3 begin
4   return X;
5 end f1;
6
7 procedure Caller is
8   X : Integer := Integer'Last;
9 begin
10  X := X + 1; -- overflow check proved.
11  X := f1(X);
12 end Caller;

```

---

<sup>1</sup>This particular case has been fixed in recent versions of GNATprove.

Clearly, an overflow must happen at line 10, resulting in an exception. The analyzer, however, proves absence of overflows in `caller`. The reason is that in the Why3 backend, the postcondition of `f1` is used as an axiom in the analysis of `caller`. The resulting theory for `caller` is an inconsistent axiom set, from which (*principle of explosion*) anything can be proven, including that false VCs are true. In such circumstances, the solver may also produce a *spurious* counterexample.

In the example above, the developer gets a warning for the inconsistent postcondition and can correct for it, thus keep barriers intact and ensure that the proofs in the caller are not influenced. However, if we change line 4 to `return x+1`, then the failing VC is now indicated in the body of `f1`, and – since the proofs build on each other – the postcondition is verified and a defect easily missed. Therefore, failing VCs within callees may also refute proofs in the caller (in contrast to execution semantics) and have to be taken into account. Indeed, the textual report of GNATprove (with flag `--assumptions`) indicates that AoRTE in `caller` depends on both the body and the postcondition of `f1`, and therefore the reports have to be studied with great care to judge the verification output. Finally, note that the same principle applies for assertions and loop invariants.

**Mistake 3: Forgetting the RTS.** Despite proven AoRTE, one procedure which rotates the frame of reference of the gyroscope measurements was sporadically triggering an exception after a floating-point multiplication. The situation was eventually captured in the debugger as follows:

```

1 | -- angle = 0.00429, vector (Z) = -2.023e-38
2 | result(Y) := Sin (angle) * vector(Z);
3 | -- result(Y) = -8.68468736e-41 => Exception

```

Variable `result` was holding a *subnormal* floating-point number, roughly speaking, an “underflow”. GNATprove models floating-point computations according to IEEE-754, which requires support for subnormals on the target processor. Our processor’s FPU indeed implements subnormals, but the RTS, part of which describes floating-point capabilities of the target processor, was incorrectly indicating the opposite<sup>2</sup>. As a result, the language-defined float validity check occasionally failed (in our case when the glider was resting level and motionless at the ground for a longer period of time). Therefore, the RTS must be carefully configured and checked manually for discrepancies, otherwise proofs can be refuted since static analysis works with an incorrect premise.

**Mistake 4: Bad Patterns.** *Saturation* may seem like an effective workaround to ensure overflows, index checks and so on cannot fail, but it usually hides bigger flaws. Consider the following example, also from the GPS protocol parser:

```

1 | subtype Lat_Type is Angle_Type range -90.0 * Degree .. 90.0 * Degree;
2 | Lat : Lat_Type := Dim_Type (toInt32 (data_rx(28..31))) * 1.0e-7 * Degree;

```

The four raw bytes in `data_rx` come from the GPS device and represent a scaled float, which could in principle carry a value exceeding the latitude range of  $[-90, 90]$  Degree. To protect against this sort of error, it is tempting to implement a function (even a generic) of the form `if x > Lat_Type'Last then x := Lat_Type'Last else...` that limits the value to the available range, and apply it to all places

<sup>2</sup>This also has been fixed in recent versions of the embedded ARM RTS.

where checks could be failing. However, we found that almost every case where saturation was applied, was masking a boundary case that needs to be addressed. In this example, we needed handling for a GPS that yields faulty values. In general, such cases usually indicate a missing software requirement.

## 4.2 Design Limitations

We now describe some cases where the current version of the SPARK 2014 *language* – not the static analysis tool – imposes limitations.

**Access types** (pointers) are forbidden in SPARK, however, low-level drivers heavily rely on them. One workaround is to hide those types in a package body that is not in SPARK, and only provide a SPARK specification. Naturally, the body cannot be verified, but at least its subprograms can be called from SPARK subprograms. Sometimes it is not possible to hide access types, in particular when packages use them as interface between each other. This is the case for our SD card driver, which is interfaced by an implementation of the FAT filesystem through access types. Both are separate packages, but the former one exports restricted types and access types which are used by the FAT package, thus requiring that wide parts of the FAT package are written in Ada instead of SPARK.

As a consequence, access types are sometimes demanding to form larger monolithic packages, here to combine SD card driver and FAT filesystem into one (possibly nested) package.

**Polymorphism** is available in SPARK, but its use is limited as a result of the access type restriction. Our message queue between flight-critical and mission-critical task was planned to hold messages of a polymorphic type. However, without access types the only option to handover messages would be to take a deep copy and store it in the queue. However, the queue itself is realized with an array and can hold only objects of the same type. This means a copy would also be an upcast to the base type. This, in turn, would loose the components specific to the derived type, and therefore render polymorphism useless. As a workaround, we used mutable variant records.

**Interfaces.** Closely related to polymorphism, we intended to implement sensors as polymorphic types. That is, specify an abstract sensor interface that must be overridden by each sensor implementation. Towards that, we declared an abstract tagged type with abstract primitive methods denoting the interface that a specific sensor must implement. However, when we override the method for a sensor implementation, such as the IMU, SPARK requires specifying the global dependencies of the overriding IMU implementation as class-wide global dependencies of the abstract method (SPARK RM 6.1.6). This happens even without an explicit `Global` aspect. As workaround, we decided to avoid polymorphism and used simple inheritance without overriding methods.

**Dimensioned Types.** Using the GNAT dimensionality checking system in SPARK, had revealed two missing features. Firstly, in the current stable version of the GNAT compiler, it is not possible to specify general operations on dimensioned types that are resolved to specific dimensions during compilation. For example, we could not write a generic time integrator function for the PID

controller that multiplies any dimensioned type with a time value and returns the corresponding unit type. Therefore, we reverted to dimensionless and unconstrained floats within the generic PID controller implementation. Secondly, it is not possible to declare vectors and matrices with mixed subtypes, which would be necessary to retain the dimensionality information throughout vector calculations (e.g., in the Kalman Filter). As a consequence, we either have split vectors into their components, or reverted to dimensionless and unconstrained floats. As a result of these workarounds, numerous overflow checks related to PID control and Kalman Filter could not be proven (which explains more than 70% of our failed floating-point VCs).

### 4.3 Solver Weaknesses

We now summarize some frequent problems introduced by the current state of the tooling.

The `'Position` attribute of a record allows evaluating the position of a component in that record. However, GNATprove has no precise information about this position, and therefore proofs building on that might fail.

Another feature that is used in driver code, are *unions*, which provide different views on the same data. GNATprove does not know about the overlay and may generate False Positives for initialization, as well as for proofs which build on the relation between views.

We had several False Positives related to possibly uninitialized variables. SPARK follows a strict data initialization policy. Every (strict) output of a sub-program must be initialized. In the current version, GNATprove only considers initialization of arrays as complete when done in a single statement. This generates warnings when an array is initialized in multiple steps, e.g., through loops, which we have suppressed.

## 5 Results

In general, verification of SPARK 2014 programs is accessible and mostly automatic. Figure 1 shows the results of our launch release. As it can be seen, we could not prove all properties during the time of this project (three months). The non-proven checks have largely been identified as “fixable”, following our design recommendations given below.

The complexity of our flight stack and verification progress are summarized in Table 1. It can be seen that our focus on the application part is reflected in the SPARK coverage that we have achieved (82% of all bodies in SPARK, and even 99% of all specifications), but also that considerably more work has to be done for the BSP (currently only verified by testing). In particular, the HAL (off-chip device drivers, bus configuration, etc.) is the largest part and thus needs a higher SPARK coverage. However, we should add that 43% of the HAL is consisting of specifications generated from CMSIS-SVD files, which do not contain any subprograms, but only definitions of peripheral addresses and record definitions



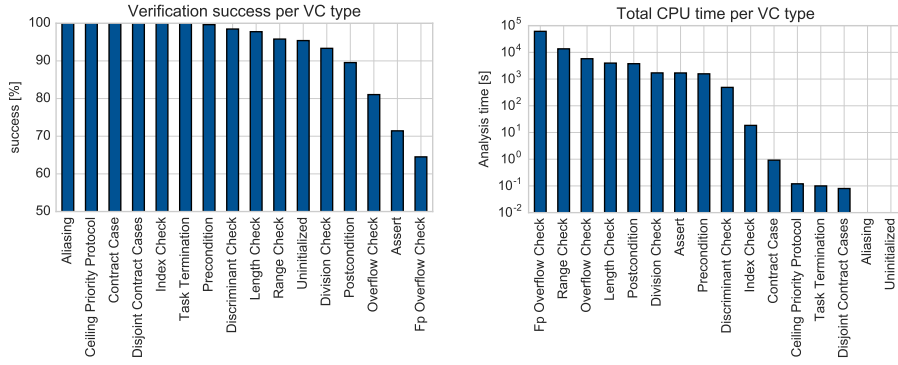


Fig. 1: Statistics on Verification Conditions (VCs) by type.

Table 1: Metrics and verification statistics of our Flight Stack.

Metric	Application	Board Support Package		All
		HAL	RTS	
lines of code (GNATmetric)	6,750	32,903	15,769	55,422
number of packages	49	100	121	270
cyclomatic complexity	2.03	2.67	2.64	2.53
SPARK body/spec	81.9/99.4 %	15.5/23.5 %	8.6/11.8 %	30.0/38.5 %
number of VCs	3,214	765	2	3,981
VCs proven	88.1 %	92.5 %	100 %	88.8 %
analysis time <sup>a</sup>	–	–	–	19 min

<sup>a</sup>Intel Xeon E5-2680 Octa-Core with 16 GB RAM, timeout=120s, steps=inf.

to access them, and therefore mostly cannot be covered in SPARK. Last but not least, a completely verified RTS would be desirable, as well.

**Floats are expensive.** Statistically, we have spent most of the analysis time (65 %) for proving absence of floating-point overflows, although these amount to only 21 % of all VCs. This is because discharging such VCs is in average one magnitude slower than discharging most other VC types. In particular, one has to allow a high step limit (roughly the number of decisions a solver may take, e.g., deciding on a literal) and a high timeout. Note that at some point an increase of either of them does not improve the result anymore.

**Multi-Threading** could be proven to follow our goals. By using the Ravenscar RTS, our goals related to deadlock, priority inversion and blocking, hold true by design. Several race conditions and non-thread-safe subprograms have been identified by GNATprove, which otherwise would have refuted task separation. To ensure that termination of low-criticality tasks cannot terminate the flight-critical task, we provided a custom implementation for GNAT's last chance handler (outside of the SPARK language and therefore not being analyzed) which reads the priority of the failing task and acts accordingly: If the priority is lower than that of the flight-critical task (i.e., the mission-critical task had an exception), then we prevent a system reset by sending the low-priority task into an infinite null loop (thus keeping it busy executing `nops`, and keeping the flight-critical task alive). If the flight-critical task is failing, then our handler allows a system reset. Multi-threading is therefore easy to implement, poses no verification problems, and can effectively separate tasks by their criticality.

**High-Level Behavioral Contracts** related to the homing functionality could be expressed and proven with the help of ghost functions, although this is beyond the main purpose of SPARK contracts. For example, we could prove the overall behavior in case of losing the GPS fix, or missing home coordinates.

## 5.1 Design Recommendations

The following constructs and strategies have been found amenable to verification:

1. Split long expressions into multiple statements → discharges more VCs.
2. Limit ranges of data types, especially floats → better analysis of overflows.
3. Avoid saturation → uncovers missing error handling and requirements.
4. Avoid interfaces → annotations for data flows break concept of abstraction.
5. Emulate polymorphic objects that must be copied with mutable variant records.
6. Separation of tasks by criticality using a custom last chance handler → abnormal termination of a low-criticality task does not cause termination of high-criticality tasks.

## 6 Related Work

Only a small number of experience reports about SPARK 2014 have been published before. A look back at (old) SPARK's history and its success, as well

as an initial picture of SPARK 2014 is given by Chapman and Schanda in [4]. We can report that the mentioned difficulties with floating-point numbers are solved in SPARK 2014, and that the goal to make verification more accessible, has been reached. A small case study with SPARK 2014 is presented in [10], but at that point multi-threading (Ravenscar) was not yet supported, and floating point numbers have been skipped in the proof. We can add to the conclusion given there, that both are easily verified in “real-world” code, although float proofs require more (computational and mental) effort. Larger case studies are summarized by Dross et al. in [5], with whom we share the opinion of minor usability issues, and that some small amount of developer training is required. Finally, SPARK 2014 with Ravenscar has recently been announced to be used in the Lunar IceCube [2] satellite, a successor of the successful CubeSat project that was implemented in SPARK 2005. It will be a message-centric software, conceptually similar to NASA’s cFE/CFS, but fully verified and striving to become an open source platform for spacecraft software. In contrast to all the above publications, this paper is not focused on the application or case studies, but pointing out typical sources of errors in SPARK programs, which a developer has to know in order to get correct verification results.

## 7 Conclusion

Although the verification of SPARK 2014 programs is very close to execution semantics and therefore mostly intuitive, we believe that developers still need some basic training to avoid common mistakes as described in this paper, which otherwise could lead to a false confidence in the software being developed. Overall, the language forces developers to address boundary cases of a system explicitly, which eventually helps understanding the system better, and usually reveals missing requirements for boundary cases. As a downside, SPARK 2014 programs are often longer than (approximately) equivalent Ada programs, since in the latter case a general exception handler can be installed to handle all pathological cases at once, without differentiating them. Furthermore, static analysis is ready to replace unit tests, but integration tests have still been found necessary.

Regarding the shortcomings of the GNAT dimensionality system, we can report that as a consequence of our experiments, a solution for generic operations on dimensioned has been found and will be part of future GNAT releases.

Our remaining criticism to SPARK 2014 and its tools is as follows: next to some minor tooling enhancements to avoid the mistakes mentioned earlier and adding some more knowledge to the analyzer, it is necessary to support object-oriented features in a better way. All in all, SPARK 2014 raises the bar for formal verification and its tools, but developers still have to be aware of limitations.

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